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BUZZARD: A NUMA-Aware In-Memory Indexing System

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ABSTRACT
With the availability of large main memory capacities, in-memory index structures have become an important component of modern data management platforms. Current research even suggests index-based query processing [2] as an alternative or supplement for traditional tuple-at-a-time processing models. However, while simple sequential scan operations can fully exploit the high bandwidth provided by main memory, indexes are mainly latency bound and spend most of their time waiting for memory accesses.

Considering current hardware trends, the problem of high memory latency is further exacerbated as modern shared-memory multiprocessors with non-uniform memory access (NUMA) become increasingly common. On those NUMA platforms, the execution time of index operations is dominated by memory access latency which increases dramatically when accessing memory on remote sockets. Therefore, good index performance can only be achieved through careful optimization of the index structure to the given topology.

BUZZARD is a NUMA-aware in-memory indexing system. Using adaptive data partitioning techniques, BUZZARD distributes a prefix-tree-based [1] index across the NUMA system and hands off incoming requests to worker threads located on each partition’s respective NUMA node. This approach reduces the number of remote memory accesses to a minimum and improves cache utilization. In addition, all indexes inside BUZZARD are only accessed by their respective owner, eliminating the need for synchronization primitives like compare-and-swap.

Figure 1 shows how BUZZARD distributes requests across NUMA nodes: (1) For each incoming request, BUZZARD uses an adaptive partition table to determine the partition the requested data belongs to. (2) Once the respective partition has been identified, the request is inserted into a thread-local intermediate request buffer. (3) Periodically, these intermediate buffers are atomically flushed to the partitions’ main request buffers, reducing contention in the main request buffers through batch insertions and hence yielding higher buffer throughput. (4) Worker threads (one per partition) located on the different NUMA nodes extract requests from the buffer of their respective partition and handle these requests by using a private index held in local main-memory. BUZZARD processes multiple requests as a batch, effectively hiding memory latency by interleaving multiple index lookups. If the workload changes, BUZZARD autonomously rebalances partitions to achieve an even load distribution.

First benchmarking results show that NUMA-aware index partitioning has a huge impact on the throughput of in-memory indexing systems running on NUMA hardware. Especially on machines with complex topologies and remote memory accesses that require multiple hops, resulting in further increased memory latency, BUZZARD speeds up index performance up to 220% compared to a single index.

Categories and Subject Descriptors
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NUMA; in-memory indexing; prefix trees

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